



Enhancing Static Analysis for Practical Bug Detection: An LLM-Integrated Approach



aka: The Hitchhiker's Guide to Program Analysis: A Journey with Large Language Models

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ABSTRACT:

- We investigate where and how **Large Language Models (LLMs)** can assist **static analysis** by asking appropriate questions. We target a specific bug-finding static analysis tool (UBITect) that produces a large number of **undecided cases** due to timeout/unknown APIs. We use LLM to handle these cases.
- Under 300 samples, we find our tool reach a precision of 50%, w/o missed bugs.
- With 1000 samples, we find 13 new bugs.

Background: UBITect

- UBITect is an analyzer and aims to find *Use-before-initialization* (UBI) bugs all over the Linux kernel.
- UBITect uses *symbolic execution* to reduce false positives, but suffers from timeout/out of memory and leaves many undecided cases.
- Our proposed tool, “LLift”, takes these undecided cases and analyze them under LLM.

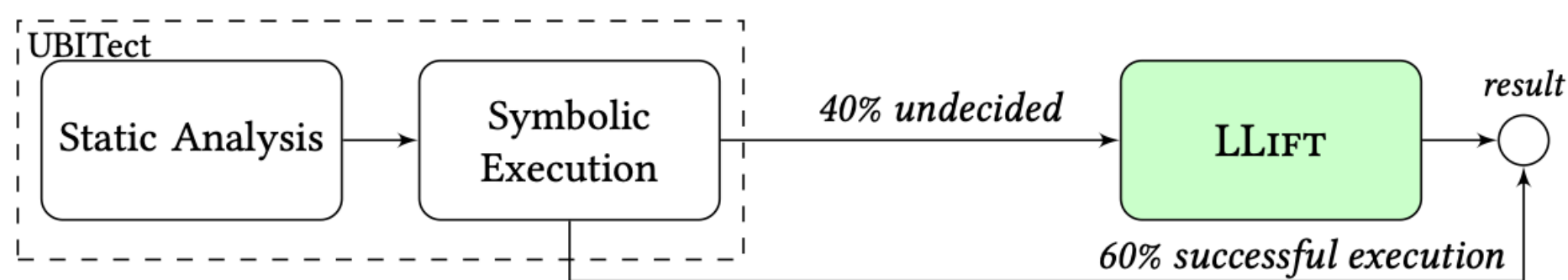


Figure 1: The high-level workflow of LLift. Take undecided cases by UBITect and determine whether these cases are real bugs.

Motivation:

```

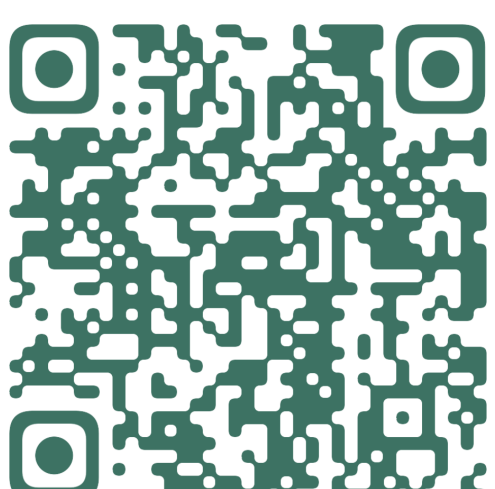
static int libcfcs_ip_str2addr(...){
    unsigned int a, b, c, d;
    if (sscanf(str, "%u.%u.%u.%u%n",
        &a, &b, &c, &d, &n) ≥ 4 && ...){
        // use of a, b, c, d
    }

    int sscanf(const char *buf, const char *fmt, ...){
        va_start(args, fmt);
        i = vsscanf(buf, fmt, args);
        va_end(args);
    }
}
  
```

Figure 2: Code snippet of sscanf and its usecase, derived from Linux kernel

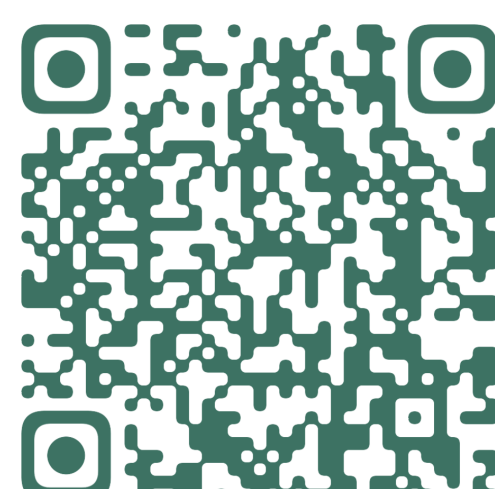
- In the sscanf case, UBITect lacks return value sensitivity (ret ≥ 4 check). It simply estimates that all parameters “*may*” be *left uninitialized*.
- LLM can resolve this case because:
 - It is familiar with function `sscanf`, knowing the meaning of *return value*.
 - With the return value ≥ 4, it can infer the `sscanf` must initialize first four parameters.

Paper Link:



<https://dl.acm.org/doi/10.1145/3649828>

Homepage:



<https://sites.google.com/view/llift-open>

Return Value Check (Post-Constraint):

Table 1: Two types of post-constraints: check before use, failure check

Check Before Use	Failure Check
Type A: <pre> if (sscanf(...) ≥ 4) { use(a, b, c, d); } </pre>	Type B: <pre> err = func(&a); if (err) { return/break/goto; } use(a) </pre>
Type A': <pre> switch(ret=func(&a)){ case some_irrelevant_case: do_something(...); break; case critical_case: use(a); } </pre>	Type B': <pre> while(func(&a)){ do_something(...); } use(a); </pre>

- More generally, not all paths are important for UBI, only “use” paths are essential.
- By focusing on return value checks (post-constraint), we can prune out many paths.

Workflow:

LLift prompts LLM to:

- Identify the “initializer”
- Extract the post-constraints
- Analyze the initializer with post-constraints

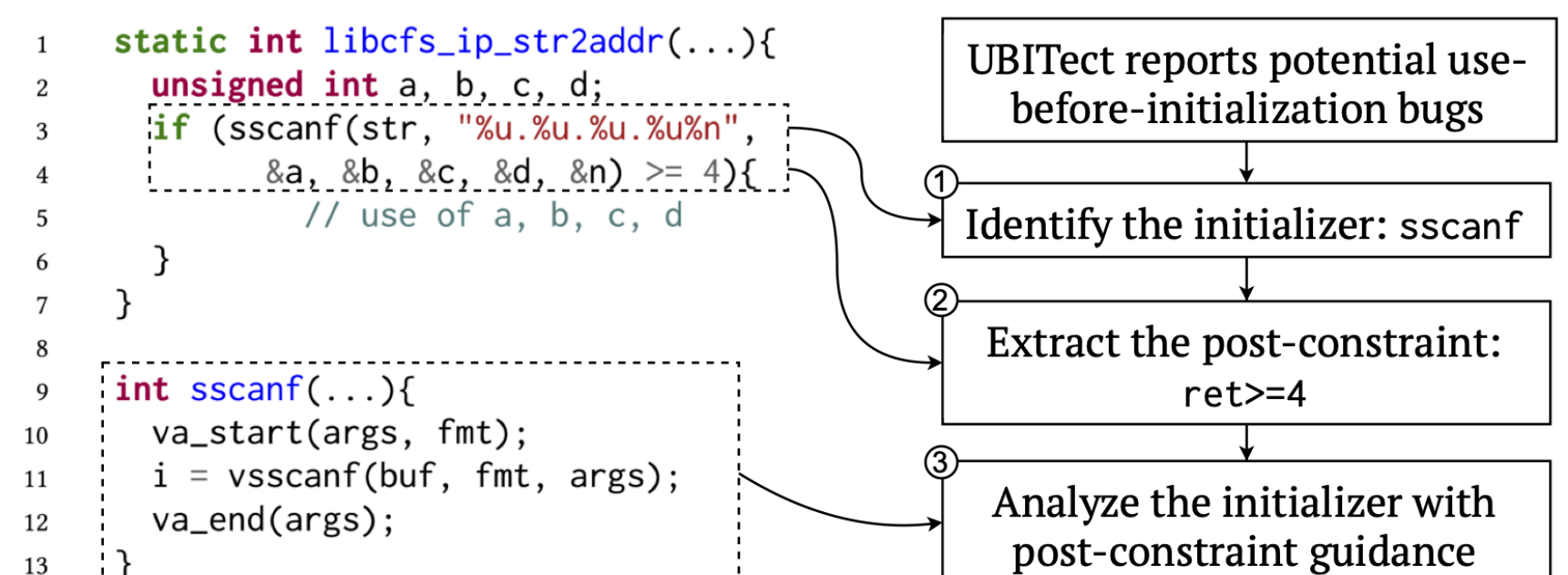


Figure 3: Example run of LLift.

Prompt Design: Progressive Prompt

- Linux kernel has a very large codebase, functions are often called across files.
- Progressive prompt:** dynamically provides function definition when LLM *needs*.

w/o progressive prompt	w/ progressive prompt
Q: summarize the function <p9_pdu_readf(..., &ecode)>, will it initialize variable <ecode>?	Q: summarize the function <p9_pdu_readf(..., &ecode)>, will it initialize variable <ecode>? if you encounter uncertainty due to a lack of function definitions, please signal your need, and I'll supply them
A: p9_pdu_readf presumably reads data from a P9 protocol data unit (PDU) ...	A: p9_pdu_readf presumably reads data from a P9 protocol data unit (PDU) ... please provide the function of p9_pdu_readf
A: conclude it in must_init ❌	Q: the definition of p9_pdu_readf is ...
	A: conclude it in may_init ✅

Figure 4: A demonstration of progressive prompt. The prompts and responses are simplified.

Result:

- 13 new bugs.
- 50% precision (on undecided cases), not found false negatives yet.
- All design components are useful: simple prompt has 12% precision and 15% recall.